

Optimal Fixed-Premise Repairs of \mathcal{EL} TBoxes

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Repairs of \mathcal{EL} TBoxes

An ontology can contain axioms that are incorrect in the underlying domain, especially if

- it was constructed from incomplete data
- or using inexact methods based on machine learning.

Such errors are detected when a reasoner generates consequences from the faulty axioms.

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Definition. Given an \mathcal{EL} TBox \mathcal{T} and a repair request \mathcal{P} (a finite set of concept inclusions), a **repair** of \mathcal{T} for \mathcal{P} is a TBox \mathcal{T}' such that

R1 $\mathcal{T}' \not\models C \sqsubseteq D$ for each unwanted consequence $C \sqsubseteq D$ in \mathcal{P} .

R2 \mathcal{T}' is somehow constructed from \mathcal{T} or somehow related to \mathcal{T} (to be specified).

A repair is **optimal** if there is no repair that is “better” w.r.t. **R2**.

State of the Art

Classical repairs:

R2 $\mathcal{T}' \subseteq \mathcal{T}$, i.e., \mathcal{T}' is constructed from \mathcal{T} by deleting axioms.

Gentle repairs:

R2 \mathcal{T}' is constructed from \mathcal{T} by weakening axioms, i.e., for each $C' \sqsubseteq D'$ in \mathcal{T}' , there is some $C \sqsubseteq D$ in \mathcal{T} such that $C \sqsubseteq D \models C' \sqsubseteq D'$.

Countermodel repairs:

R2 There is a countermodel \mathcal{J} such that $\mathcal{T}' \models C \sqsubseteq D$ iff $\mathcal{T} \models C \sqsubseteq D$ and $\mathcal{J} \models C \sqsubseteq D$.

R. Reiter: **A theory of diagnosis from first principles**. Artif. Intell. (1987)

F. Baader, F. Kriegel, A. Nuradiansyah, R. Peñaloza: **Making repairs in description logics more gentle**. KR 2018

F. Kriegel: **Constructing and extending description logic ontologies using methods of formal concept analysis**. Doctoral thesis (2019)

Two New Types of Repairs

Inspired by the gentle repairs w.r.t. \succ^{sub} as well as by the countermodel repairs, and in order to tackle their problems, a novel type of repairs is introduced.

Generalized-conclusion repairs (GC-repairs):

R2 For each $C' \sqsubseteq D'$ in \mathcal{T}' , there is $C \sqsubseteq D$ in \mathcal{T} such that $C = C'$ and $\emptyset \models D \sqsubseteq D'$.

Theorem. For each TBox \mathcal{T} and each repair request \mathcal{P} ,

- the set of all optimal GC-repairs can be computed in exponential time,
- and every GC-repair is entailed by an optimal one.

Two New Types of Repairs

We introduce yet another type of repairs which can retain more consequences than GC-repairs.

Fixed-premise repairs (FP-repairs):

R2 $\mathcal{T} \models \mathcal{T}'$ and for each $C' \sqsubseteq D'$ in \mathcal{T}' , there is $C \sqsubseteq D$ in \mathcal{T} such that $C = C'$.

Theorem. For each TBox \mathcal{T} and each repair request \mathcal{P} ,

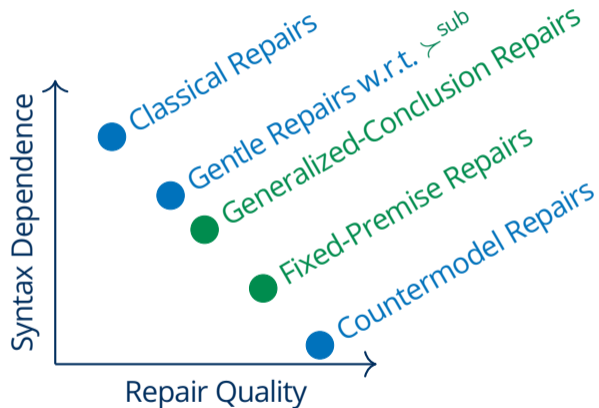
- the set of all optimal FP-repairs can be computed in exponential time,
- and every FP-repair is entailed by an optimal one.

Computing Repairs

Canonical GC- and FP-repairs can be constructed in almost the same way:

- 1 Choose a polynomial-size **repair seed** \mathcal{S} .
- 2 Construct the **induced countermodel** $\mathcal{J}_{\mathcal{S}}$.
- 3 For each concept inclusion $C \sqsubseteq D$ in \mathcal{T} ,
 - GC: generalize D such that the resulting concept inclusion holds in $\mathcal{J}_{\mathcal{S}}$.
 - FP: replace D with the most specific concept E for which the concept inclusion $C \sqsubseteq E$ holds in $\mathcal{J}_{\mathcal{S}}$.

Comparison of \mathcal{EL} TBox Repair Approaches



An Example with Bikes

Example. We consider the \mathcal{EL} TBox

$$\mathcal{T}_{\text{Bike}} := \left\{ \begin{array}{l} \text{MountainBike} \sqsubseteq \text{Bike}, \\ \text{Bike} \sqsubseteq \exists \text{hasPart.SuspensionFork} \sqcap \exists \text{isSuitableFor.OffRoadCycling}, \\ \text{SuspensionFork} \sqsubseteq \text{Fork}, \\ \text{OffRoadCycling} \sqsubseteq \text{Cycling} \end{array} \right\}$$

The TBox $\mathcal{T}_{\text{Bike}}$ entails two faulty consequences

- 1 $\text{Bike} \sqsubseteq \exists \text{hasPart.SuspensionFork}$
 - 2 $\text{Bike} \sqsubseteq \exists \text{isSuitableFor.OffRoadCycling}$
- } repair request $\mathcal{P}_{\text{Bike}}$

An Example with Bikes: Classical Repair

Example. A classical repair of $\mathcal{T}_{\text{Bike}}$ for $\mathcal{P}_{\text{Bike}}$ is

$$\left\{ \begin{array}{l} \text{MountainBike} \sqsubseteq \text{Bike}, \\ \text{Bike} \sqsubseteq \exists \text{hasPart.SuspensionFork} \sqcap \exists \text{isSuitableFor.OffRoadCycling}, \\ \text{SuspensionFork} \sqsubseteq \text{Fork}, \\ \text{OffRoadCycling} \sqsubseteq \text{Cycling} \end{array} \right\}$$

An Example with Bikes: Generalized-Conclusion Repair

Example. An optimal GC-repair of $\mathcal{T}_{\text{Bike}}$ for $\mathcal{P}_{\text{Bike}}$ is

$$\left\{ \begin{array}{l} \text{MountainBike} \sqsubseteq \text{Bike}, \\ \text{Bike} \sqsubseteq \exists \text{hasPart.SuspensionFork} \sqcap \exists \text{isSuitableFor.OffRoadCycling}, \\ \quad \exists \text{hasPart.T} \sqcap \exists \text{isSuitableFor.T}, \\ \text{SuspensionFork} \sqsubseteq \text{Fork}, \\ \text{OffRoadCycling} \sqsubseteq \text{Cycling} \end{array} \right\}$$

It is also a gentle repair of $\mathcal{T}_{\text{Bike}}$ for $\mathcal{P}_{\text{Bike}}$ w.r.t. the weakening relation \succ^{sub} .

An Example with Bikes: Fixed-Premise Repair

Example. An optimal FP-repair of $\mathcal{T}_{\text{Bike}}$ for $\mathcal{P}_{\text{Bike}}$ is

$$\left\{ \begin{array}{l} \text{MountainBike} \sqsubseteq \text{Bike}, \\ \text{Bike} \sqcap \exists \text{hasPart.SuspensionFork} \sqcap \exists \text{isSuitableFor.OffRoadCycling}, \\ \text{Bike} \sqsubseteq \exists \text{hasPart.SuspensionFork} \sqcap \exists \text{isSuitableFor.OffRoadCycling}, \\ \quad \exists \text{hasPart.Fork} \sqcap \exists \text{isSuitableFor.Cycling}, \\ \text{SuspensionFork} \sqsubseteq \text{Fork}, \\ \text{OffRoadCycling} \sqsubseteq \text{Cycling} \end{array} \right\}$$

Do you have questions or comments?